

## Problem Statement

- When data is stored in cloud or untrusted remote server, it is very challenging to share that data securely if multiple groups of user exist.
- Designing a data sharing scheme in such a scenario needs to achieve following goals:
  - Scheme should be scalable with number of user
  - Member leaving or new member joining cost should be minimal
  - Ensures Forward and Backward Secrecy
  - Group level data isolation
  - Cross-group data sharing
  - Store and share data securely within group members using existing untrusted public cloud
  - Fine grained access control in shared data
  - Prevent collusion attack

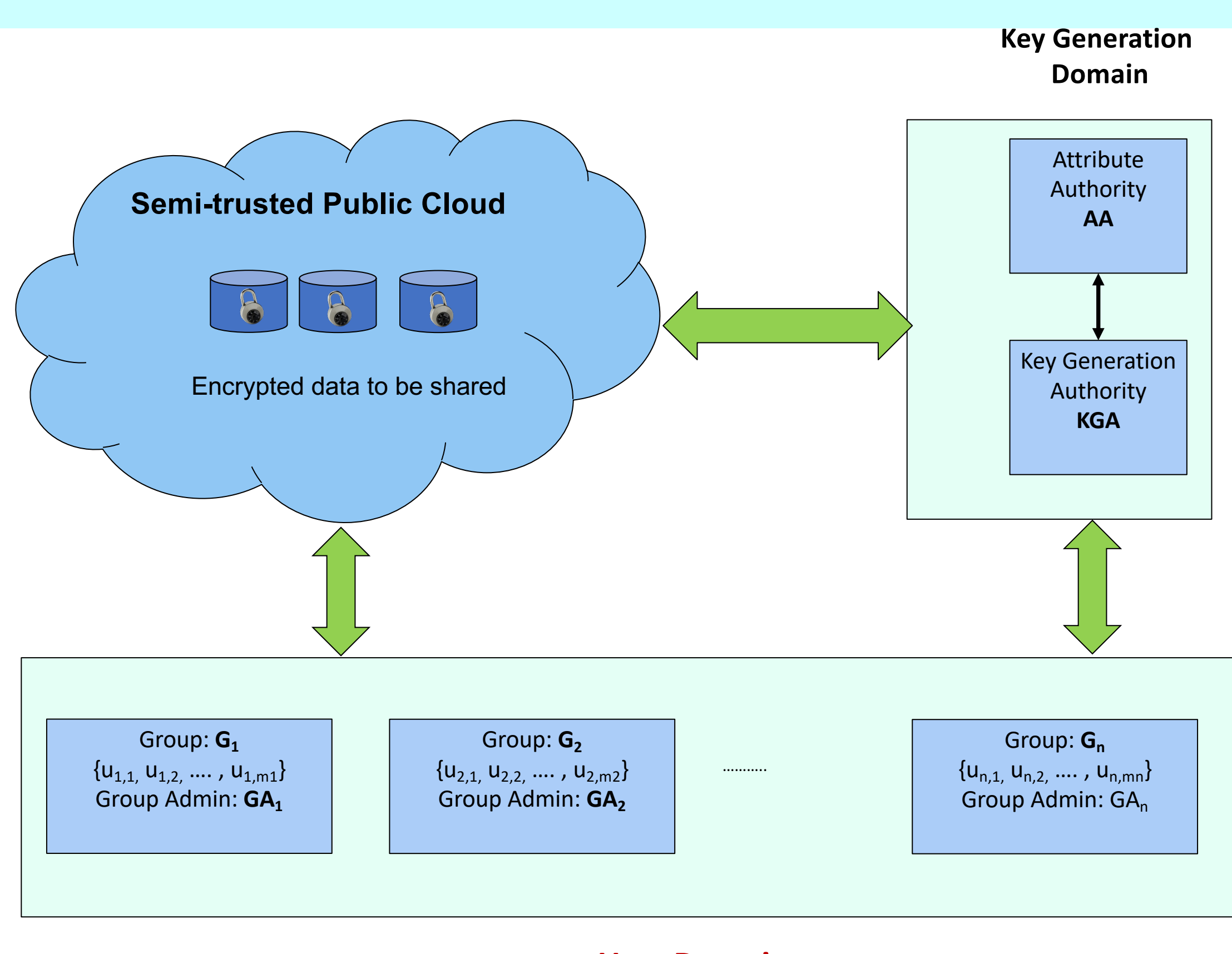
## Challenges

- Handle membership change event without affecting keys of currently active users
- Enable on demand cross-group data sharing at file level granularity without affecting all files
- Preventing key-escrow problem

## Threat Model

- Public Cloud will try to learn plaintext from stored cypher-text
- Member of one group with same attributes should not be able decrypt data of other group
- A compromised user other than the file owner will try to modify access policy of the file
- Multiple users may collude with each other and try to decrypt cypher text that can not be decrypted individually
- Revoked user may collude with cloud to decrypt data
- Assumption:
  - Cloud is semi-trusted that means it follows the protocols
  - User of one group does not share his TGDH key tree secret key with members of other groups

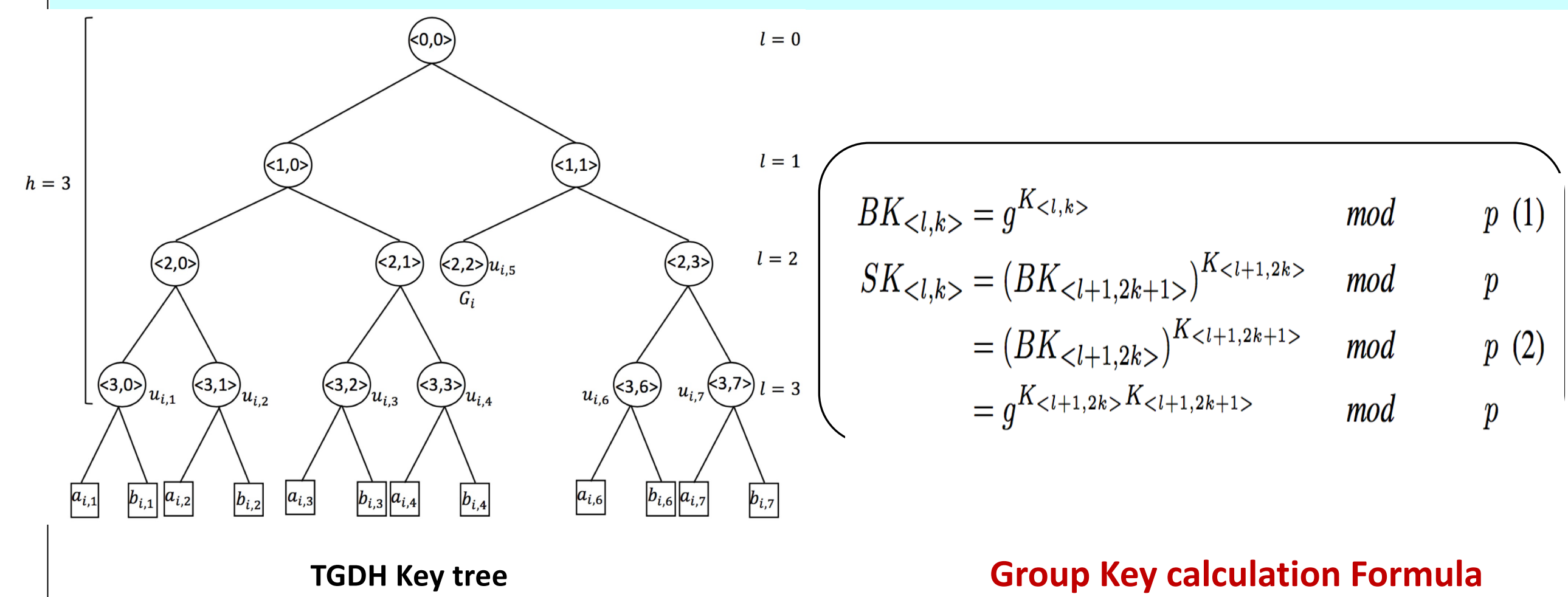
## System Architecture



Scheme	Security assumption	Model	Outsourced Decryption	Verifiability	Revocation	Unlimited joining	Multi-group	Collusion resistant
DASS [8]	Decisional PB-DHE	Standard	X	X	✓	X	X	X
Hur-I [13]	Generic Group	RO	X	X	✓	X	X	X
Hur-II [7]	Generic Group	RO	X	X	✓	X	X	X
PIRATTE [6]	Generic Group	RO	X	X	✓	✓	X	X
VO-ABE [17]	Decisional $q$ -PBDHE	Standard	✓	✓	X	X	X	N/A
CryptCloud+[5]	$l$ -SDH	Standard	X	X	✓	X	X	✓
Flexible [12]	Generic Group	RO	✓	X	✓	✓	X	✓
UserCol [14]	Generic Group	RO	X	X	✓	X	X	✓
Ours	CDH	RO	✓	✓	✓	✓	✓	✓

Comparison with related schemes in terms of security and functionality.

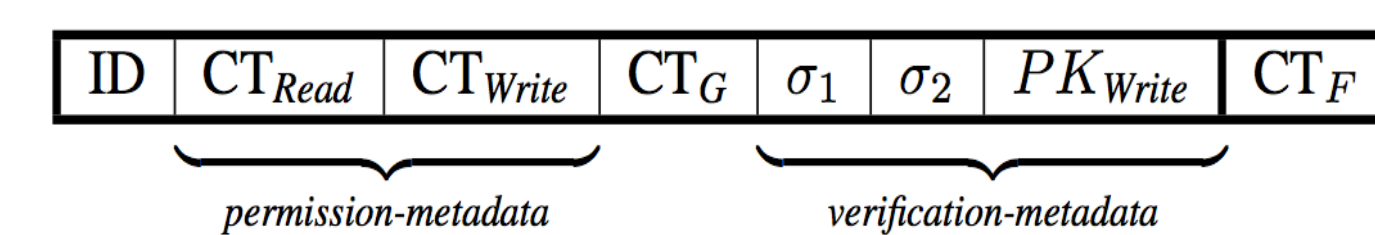
## Group Key Calculation



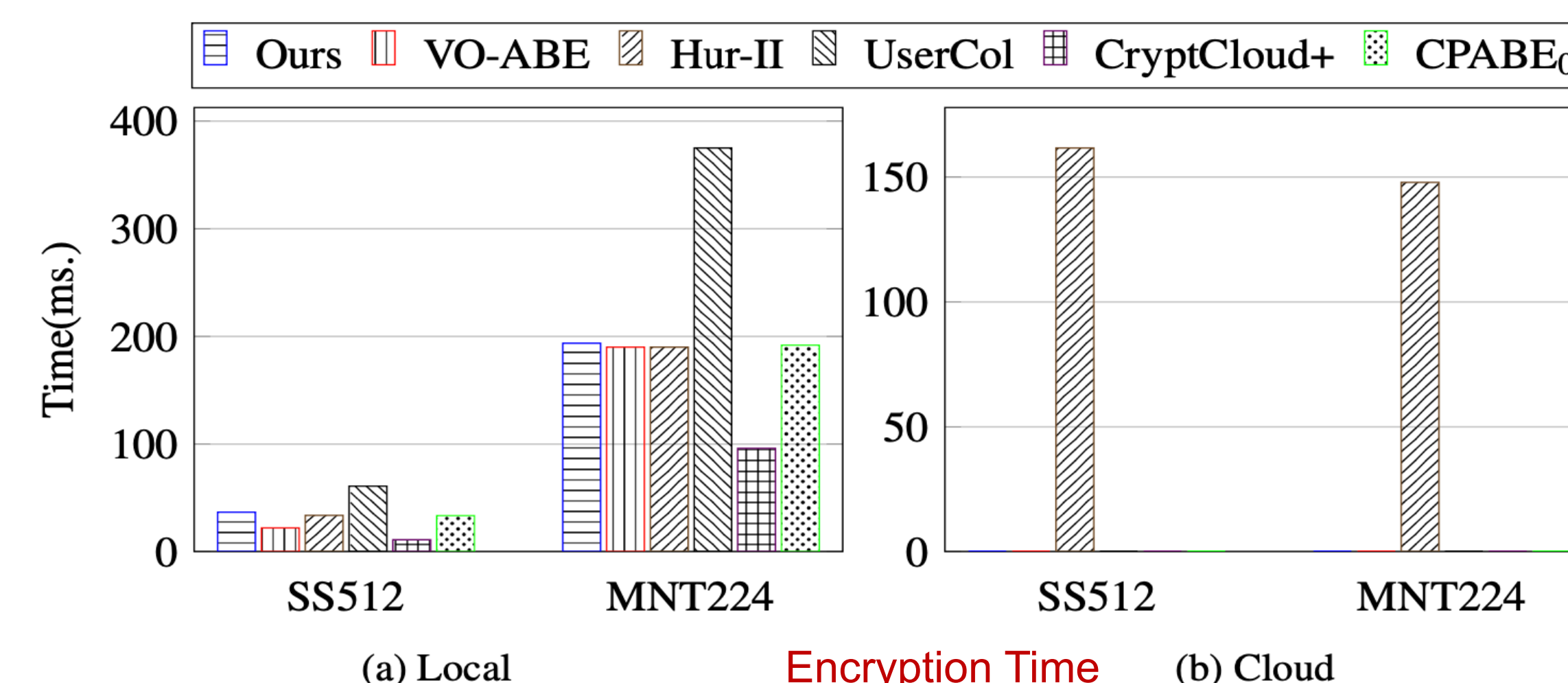
Round	level $l$	Secret key $K_{<l,k>}$	Blinded key $BK_{<l,k>}$
1	2	$K_{<2,0>} = (BK_{<3,0>})^{K_{<3,1>}} = g^{K_{<3,0>}K_{<3,1>}}$	$BK_{<2,0>} = g^{K_{<2,0>}}$
		$K_{<2,1>} = (BK_{<3,2>})^{K_{<3,3>}} = g^{K_{<3,2>}K_{<3,3>}}$	$BK_{<2,1>} = g^{K_{<2,1>}}$
		$K_{<2,2>} = s_i$	$BK_{<2,2>} = g^{s_i}$
		$K_{<2,3>} = (BK_{<3,6>})^{K_{<3,7>}} = g^{K_{<3,6>}K_{<3,7>}}$	$BK_{<2,3>} = g^{K_{<2,3>}}$
2	1	$K_{<1,0>} = (BK_{<2,0>})^{K_{<2,1>}} = g^{K_{<2,0>}K_{<2,1>}}$	$BK_{<1,0>} = g^{K_{<1,0>}}$
		$K_{<1,1>} = (BK_{<2,2>})^{K_{<2,3>}} = g^{K_{<2,2>}K_{<2,3>}}$	$BK_{<1,1>} = g^{K_{<1,1>}}$
3	0	$K_{<0,0>} = (BK_{<1,0>})^{K_{<1,1>}} = g^{K_{<1,0>}K_{<1,1>}}$	$BK_{<0,0>} = g^{K_{<0,0>}}$

## Technical Details

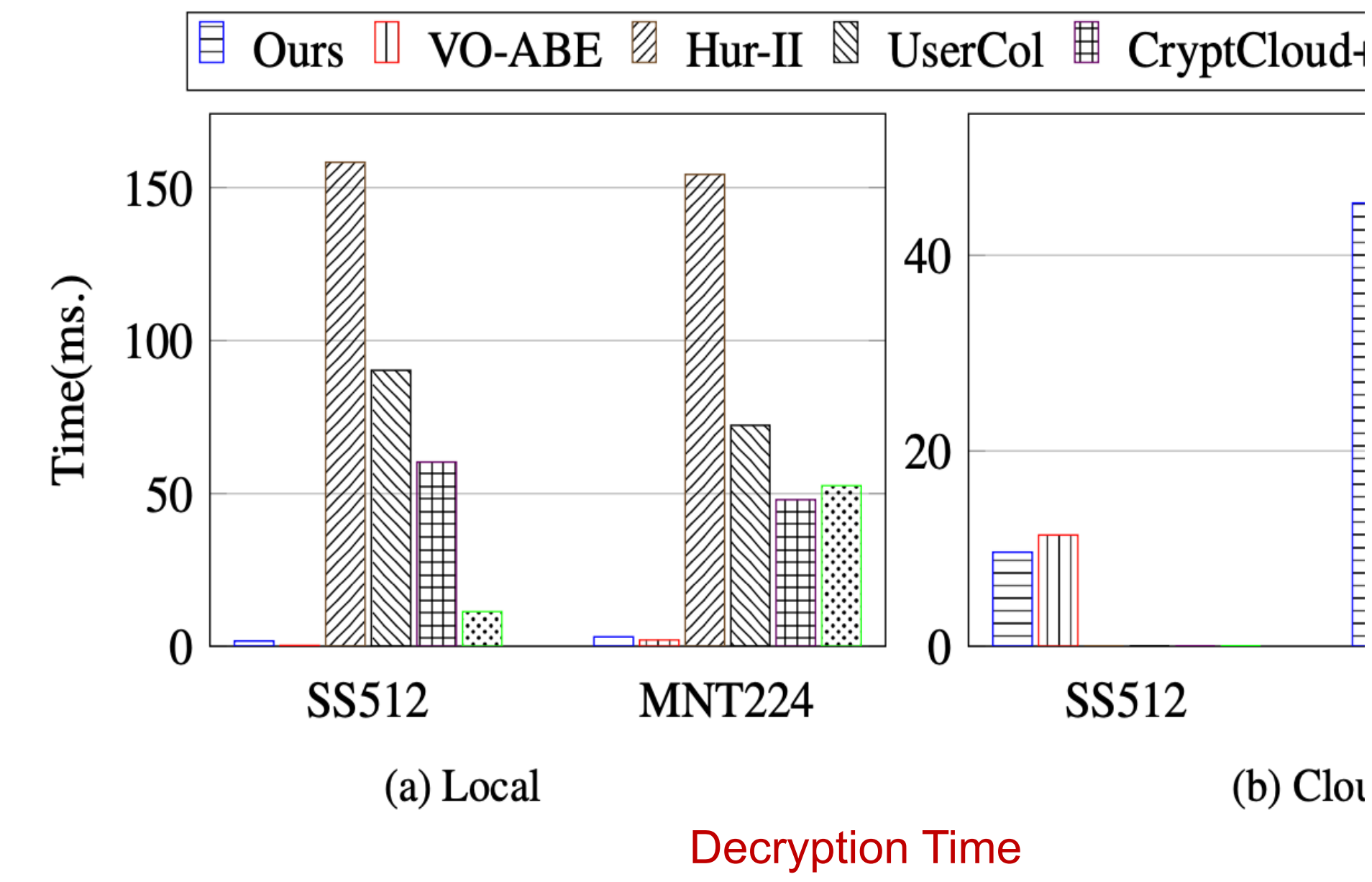
- Data owner will encrypt a file using symmetric encryption key  $K$  as  $CT_F$
- Then  $K$  is encrypted using our key escrow-free CP-ABE scheme as  $CT_{Read}$
- A file signature key  $K_{Write}$  is chosen and it is also encrypted using CP-ABE scheme as  $CT_{Write}$
- Possession of  $K$  gives one read permission on file while  $K_{Write}$  gives one write permission.
- Owner creates signature on  $\langle ID, CT_{Read}, CT_{Write}, PK_{Write} \rangle$  with his signing key
- Owner also creates a signature on encrypted file  $CT_F$
- Group denominator secret is encrypted with the current TGDH public key
- Owner then send all the information as the ciphertext  $CT_{Full}$  to the cloud and cloud stores the encrypted file as following format:



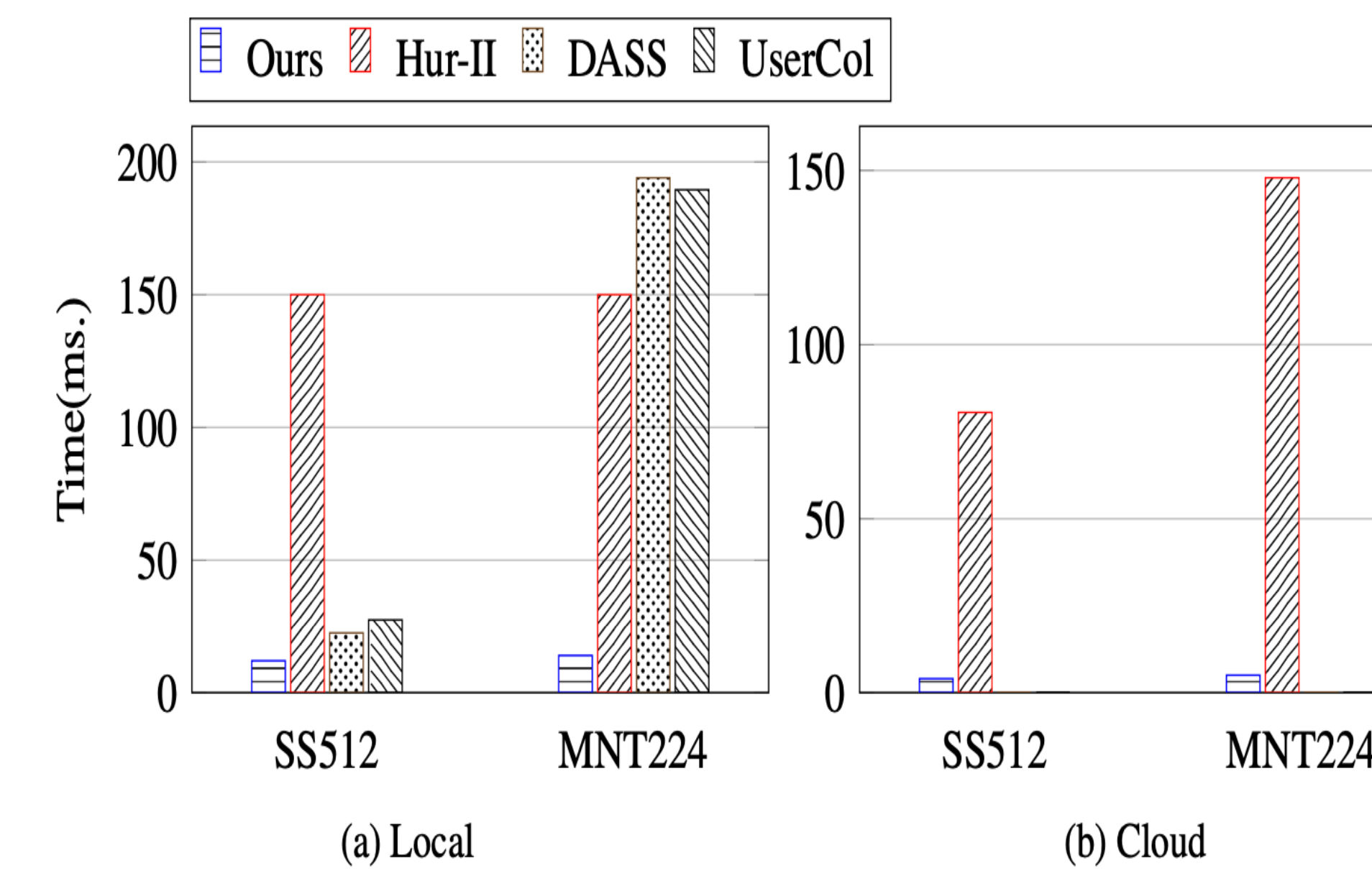
- Later, any legitimate user can download encrypted files from cloud and decrypt it if he has proper access right.



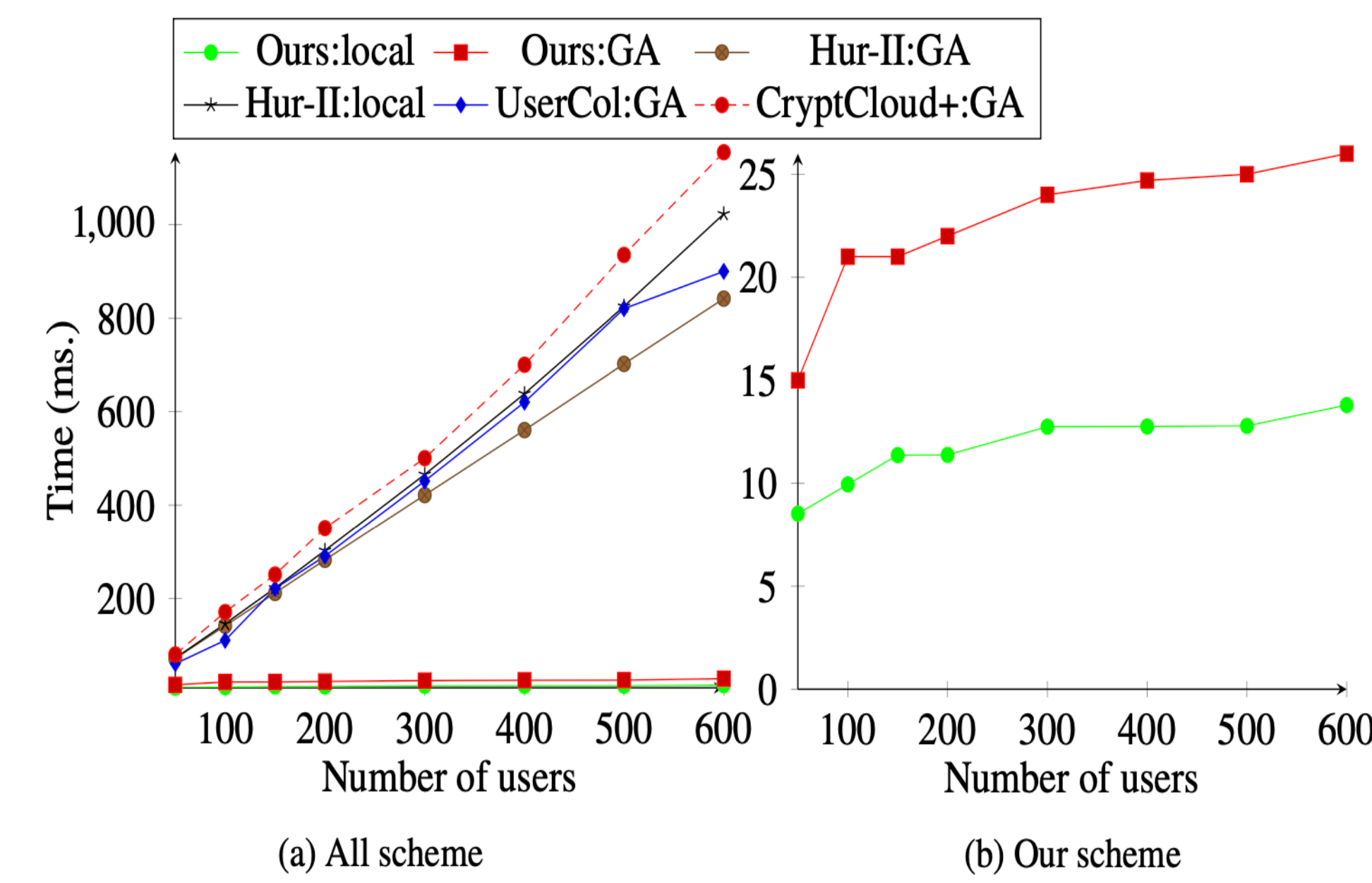
## Experiment Results



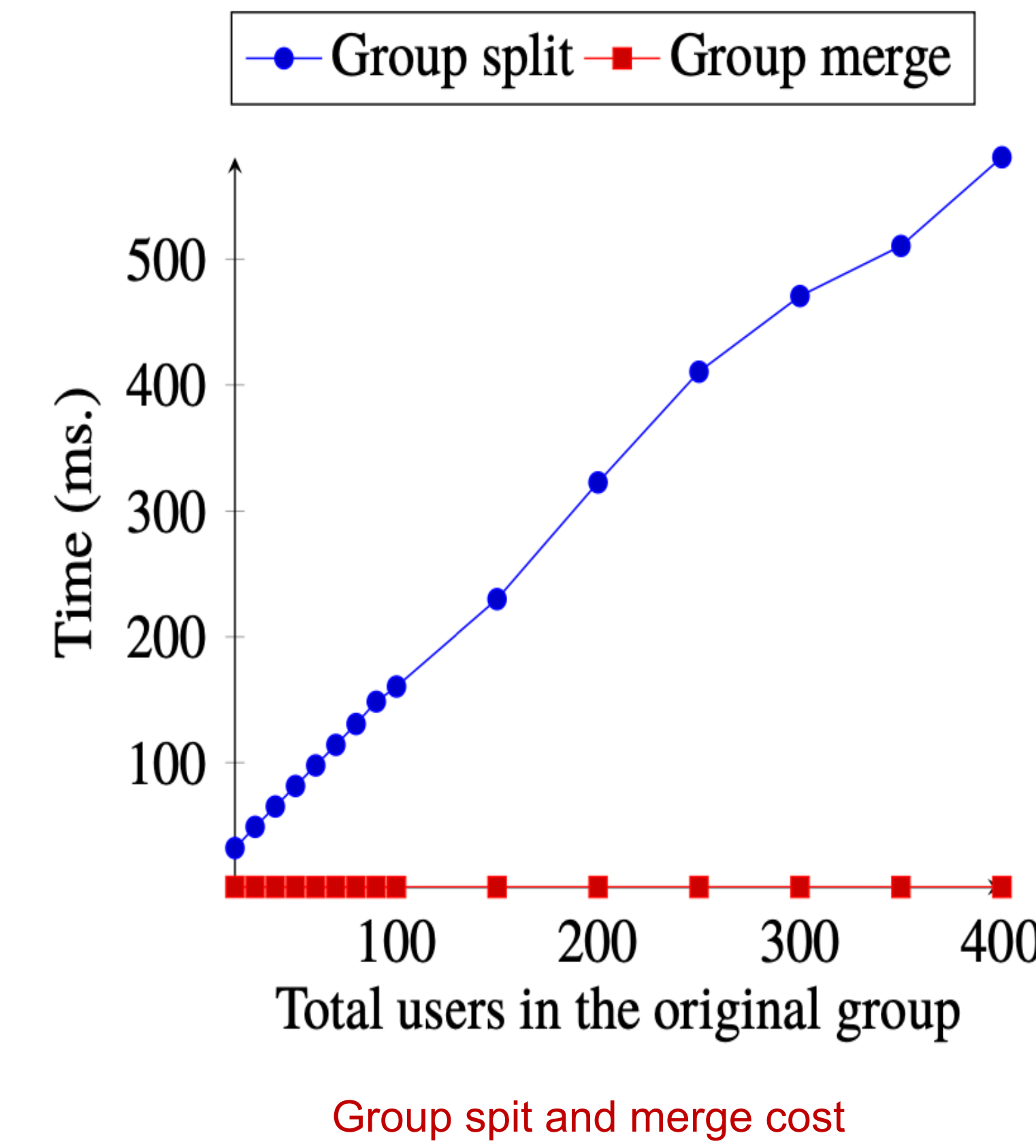
Decryption Time



Re-encryption



Re-keying time



Group split and merge cost

Scheme	Ciphertext size	Secret key size	Public key size
DASS [8]	$(2a + 1) G_1  +  G_T  +  C $	$(b + 1) G_1  + (\log m) K $	$(u + 2) G_1  +  G_T $
Hur-I [13]	$(2a + 1) G_1  +  G_T  +  C $	$(2b + 1) G_1  + (\log m) K $	$2 G_1  +  G_T $
Hur-II [7]	$(2a + 1) G_1  +  G_T  +  C $	$2(b + 1) G_1 $	$3 G_1  +  G_T $
VO-ABE [17]	$(2a + 1) G_1  +  G_T  +  C  + l_2$	$(b + 3) G_0  +  p $	$(u + 2) G_1  +  G_T $
PIRATTE [6]	$(a + 1) G_1  + a G_2  +  G_T  +  C $	$2b G_1  + (b + 1) G_2  + 2 p $	$2 G_1  +  G_2  +  G_T $
CryptCloud+ [5]	$(2a + 5) G_1  +  G_T  +  C $	$(b + 4 + 2\log m) G_1 $	$(u + 6) G_1  + 3 p $
Flexible [12]	$(2a + 6) G_1  +  G_T  + 2 p  +  C $	$(b + 4) G_1  + 2 p $	$3 G_1  + 2 G_T  +  p $
UserCol [14]	$(4a + ra + 1) G_1  +  G_T  +  C $	$4b G_1  +  G_T $	$2(u + 3) G_1  + 2 G_T  + (2m - 1) p $
Ours	$(2a + 1) G_1  +  G_T  +  C  + l_2$	$2b G_1  + 2 p $	$2 G_1  +  G_T $

Comparison of storage and communication efficiency with other schemes

Scheme	Encryption	Decryption	
		local	cloud
DASS [8]	$(3a + 1)C_1 + C_T$	$(s + 1)P + s(C_1 + C_T)$	N/A
Hur-I [13]	$(3a + 1)C_1 + C_T$	$(2s + 1)P + C_1 + C_T \log a$	N/A
Hur-II [7]	$(3a + 2m + 3)C_1 + C_T$	$(3s + 1)P + C_T \log a + (m + 1)sC_1$	N/A
VO-ABE [17]	$(3a + 1)C_1 + C_T$	$C_T$	$(2s + 1)P + aC_T$
PIRATTE [6]	$(a + 1)C_1 + C_T + aC_2$	$(s + \log a)C_T + (3s + 1)P$	$aC_2$
CryptCloud+ [5]	$(a + 5)C_1 + C_T$	$2C_1 + sC_T + (2s + 5)P$	N/A
Flexible [12]	$2(a + 3)C_1 + 2C_T$	$4C_T$	$(2s + 4)P + C_T \log a$
UserCol [14]	$(3a + ra + 1)C_1 + C_T + P$	$(2s - 1)C_T + (3s + 1)P$	N/A
Ours	$2(a + 1)C_1 + C_T + P$	$2C_T$	$(2s + 1)P + C_1 + C_T \log a$

Comparison with other schemes in terms of computation cost.

Scheme	Key update		Re-encryption	
	user	Cloud/GA	Owner	Cloud/GA
DASS [8]	$bc_1$	0	$(3a + 1)C_1 + C_T$	0
Hur-I [13]	$bc_1$	0	0	$(3a + 1)C_1 + C_T$
Hur-II [7]	$bmC_1$	$2(m + 1)C_1$	0	$(3a + 1)C_1 + C_T$
PIRATTE [6]	0	$amC_2$	$(a + 1)C_1 + C_T + aC_2$	0
CryptCloud+ [5]	0	$3mC_1$	N/A	N/A
Flexible [12]	$(b + 1)C_1$	$2mC_1 + P$	0	$C_1 + P$
UserCol [14]	0	$(2m - 1)C_1$	0	$(3a + 2)C_1 + C_T$
Ours	$C_1 \log m$	$(2\log m + 3)C_1$	$C_1 + C_T + 2P$	0

Cost of group dynamic change

## CONCLUSIONS AND FUTURE WORK

- We proposed a directly revocable ABE scheme called ReVO-ABE using our proposed data structure called e-TGDH tree.
- ReVO-ABE does not put any cap on the number of user revocation or joining.
- A federated cloud architecture (using two clouds) and a novel key binding technique to prevent collusion attacks and achieve revocation under the assumption that at least one of the two clouds acts honestly
- A multi-group secure data sharing scheme called DMG-SDS to demonstrate that our ABE scheme supports a multi-group setting.
- We have only considered static access policy in this work; it will be interesting to see how it affects our system if dynamic access policy change is allowed.