CANSentry: Securing CAN-Based CPS against Denial and Spoofing Attacks

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- The most common communication protocol for automotive and industrial applications
- On-Board Diagnostics (OBD-II) is mandated to be deployed in all cars for emission control with CAN only
- It allows data transmission in hostile environments
- Due to its bus topology, it greatly reduces vehicles' cost and weight

Without CAN

With CAN



https://www.quora.com/Why-is-CAN-protocol-preferred-to-be-used-in-automotive-application



	1 2 3	4 5	6 7 8
	OBD-II	Port	
	9 10 11	12 13	14 15 16
PIN	DESCRIPTION	PIN	DESCRIPTION
1	Vendor Option	9	Vendor Option
2	J1850 Bus +	10	j1850 BUS
3	Vendor Option	11	Vendor Option
4	Chassis Ground	12	Vendor Option
5	Signal Ground	13	Vendor Option
6	CAN (J-2234) High	14	CAN (J-2234) Low
			T00 0141 0 T
7	ISO 9141-2 K-Line	15	150 9141-2 Low

OBD-II Port Gateway Cyber-Physical Cyber-Physical **B** Physical

- Conventional peer-to-peer communication paradigm
 - Hey! I'm A, this is a message to C



- In CAN
 - Hey! This is a message with ID 10 to everyone



- Format of CAN frames
 - Each frame is identified by its arbitration ID
 - The frame with the lowest ID wins the arbitration and dominates the bus
 - Different types of frames use different ID
 - Ideally, IDs should be used uniquely across ECUs













Existing Attacks

Denial attacks

[S&P'10, CCS'16, DIMVA'17, ARES'17, ESORICS'17]

- Bus Denial (BD): 0x0 ID, dominant bits via Test Mode exploitation or custom ECU
- ECU Denial (ED) : CAN Controller abuse or bypass & Error Handling abuse
- Arbitration (AD) : injection of high priority IDs, dominant bit, or fake partial frames

Spoofing attacks

[S&P'10, DefCon'13, arXiv preprint'19, BlackHat'15]

- An attacker sends any CAN ID of her choice to spoof other ECUs
- ECUs could be compromised through a remote channel, and CAN frames are sent to unlock doors, stall the engine, or control the steering wheel

Existing Controls

• Controls

- Node identification and IDS [USENIX Sec'16, CCS'17, CCS'18, TIFS'18, ACSAC'19]
- CAN-ID Obfuscation [escar'15, SCAV'17, TODAES'17, Access'19]
- Counterattacking [VTC'12, escar'14, SafeComp'18]
- Authentication [DATE'09, DATE'13, ICCPS'13]
- Firewalls [Micro'18]

Problems with existing controls

- Many require major software & hardware modifications
- Changes to the protocol may require ALL ECUs to be updated
- May introduce overheads for key management and crypto operations
- Cannot defend against abused or compromised CAN controllers
- None can handle attacks based on incomplete frames

Existing Controls

Control	Feat	ures			Effectiveness against attacks									
	Inj.	Aper.	RT	Cost	BD1	BD2	BD3	ED1	ED2	ED3	AD1	AD2	AD3	Spoof
Anomaly-based IDS	×	\checkmark	X	\checkmark	D	D	D	D	D	D	D	D	D	D
Voltage-based IDS	×	\checkmark	X	X	D	_	_	_	D	-	D	-	_	D
Time-based IDS	X	×	X	\checkmark	D	_	_	_	D	_	D	_	_	D
ID Obfuscation	X	\checkmark	X	\checkmark	_	_	_	P	P	P	P	P	P	P
Counterattacking	×	\checkmark	\checkmark	\checkmark	P	_	_	_	P	_	P	_	_	P
Authentication	X	\checkmark	X	X	P	_	_	_	D	_	D	-	_	P
Application-level Firewall	×	\checkmark	X	X	P	_	_	_	D	_	P	_	_	Р

Features: Inj.: preventing injection of incomplete frames or random bits, Aper.: handling aperiodic attacks, RT: real-time defense; Cost: low cost. Effectiveness: D: Detect, P: Prevent, -: No protection

Threat Model

 Attackers have remote access (via wireless access points) or brief physical access (via OBD-II port) to the CAN bus

• The CAN Abuser

- Has complete control over ECU's software but not hardware
- Abuse arbitration and error handling mechanisms to achieve malicious goals

• The Skipper

- Skips CAN controller to directly access CAN bus
 - Uses a custom MCU directly connected to the bus
 - Manipulate CAN controller's GPIO pins to directly access to the bus
- Attacker does not comply with CAN standards

Threat Model

- Denial attacks: disable certain functionalities in a target ECU or bus
 - ECU is shutdown (bus-off state)
 - Bus is occupied
 - Specific CAN ID cannot win arbitration
- **Spoofing attacks**: transmit an ID belonging to another ECU
 - Receiving ECUs are spoofed resulting in:
 - Disabling brakes
 - Taking control of the steering wheel
 - False data injection

A Novel Stealthy Arbitration Denial Attack

- New objectives: Selective, Stealthy, & Practical
- Overview of the attack
 - Passively monitors the bus to detect a targeted ID in the arbitration phase
 - Overwrites the last recessive bit in the target ID to win arbitration
 - Completes the transmission with a fake frame
- Challenges
 - Existing tools only deal with complete CAN frames
 - High degree of precision is needed
 - Unexpected delays, premature injection, or malformed frames may cause incomplete frames resulting in bus errors

A Novel Stealthy Arbitration Denial Attack



The Skipper



A Novel Stealthy Arbitration Denial Attack

- Stealthy the new attack does not incur any error
- Selective the new attack only affects the targeted ECU or CAN ID



CANSentry: Overview

- A novel firewall sitting between any high-risk ECU and the bus
 - High-risk ECU: an ECU with remote access (entertainment system, Bluetooth) or open hardware access (OBD-II)
- Monitors incoming traffic from the ECU
- Ensures the consistency between the CAN Bus state and the ECU state
 - E.g., when another ECU wins arbitration and transmits data, the protected ECU could only receive. It cannot interrupt the BUS traffic.
- Uses firewall rules to block illegal traffic
- Low-cost and highly efficient implementation

CANSentry: Architecture



CANSentry: States

• States and states transitions of CAN bus and nodes



CANSentry: State Transition Rules

• Main principle: The fundamental principle of the firewall is to ensure that at any time high-risk nodes on the external bus operate in a state consistent with the state of the internal bus.

CAN _{INT} State	Consistent State in CAN _{EXT}
IDLE	IDLE/RECIEVE, ARBITRATION, and TRANSMIT
ARBITRATION	IDLE/RECIEVE and ARBITRATION
TRANSMIT or ERROR FLAG	IDLE/RECEIVE

Example: R_1 : When the internal bus is in either TRANSMIT_{int} or ERRORFLAG state, the firewall always forwards the traffic from CAN_{int} to CAN_{ext} and blocks the traffic from external to internal, regardless of high-risk node's state.

CANSentry: CAN ID filtering in arbitration

- When the internal bus is in ARBITRATION state, the firewall forwards traffic that has a CAN ID in the **arbitration whitelist** and conforms to CAN specifications from CAN_{ext} to CAN_{int}.
- Prevent spoofing attacks
- Uses automata for efficient CAN ID matching
 - Example: this automata allows four CAN IDs: 0x123, 0x456, 0x789, 0x7AB

CANSentry: Implementation



CANSentry: Evaluation



Arbitration Denial Attack

Stealthy Arbitration Denial Attack

Security Analysis

- CANSentry nodes are deployed in a physically secure environment
 - Makes it difficult for an adversary to bypass or alter
- CANSentry only has two network interfaces CAN_{ext} and CAN_{int}
 - The limited communication channel and the simplicity of CAN makes it impractical to compromise the operations of the firewall from CAN_{ext} .
- The simplicity of the firewall makes it unlikely to have significant software faults

Security Analysis

-	Control	Feat	ures			Effectiveness against attacks										
-		Inj.	Aper.	\mathbf{RT}	Cost	BD1	BD2	BD3	ED1	ED2	ED3	AD1	AD2	AD3	Spoof	
Anor	maly-based IDS	×	\checkmark	X	\checkmark	D	D	D	D	D	D	D	D	D	D	
Volt	tage-based IDS	X	\checkmark	X	×	D	-	_	-	D	—	D	-	_	D	
Т	ime-based IDS	X	×	X	\checkmark	D	_	_	_	D	_	D	_	_	D	
I	D Obfuscation	X	\checkmark	X	\checkmark	_	_	_	P	P	P	P	P	P	P	
Со	ounterattacking	X	\checkmark	\checkmark	\checkmark	P	_	_	_	P	_	P	_	_	Р	
A	Authentication	X	\checkmark	X	X	P	_	_	_	D	_	D	_	_	Р	
Application	n-level Firewall	X	\checkmark	X	X	P	_	_	-	D	—	P	-	-	Р	
	CANSentry	\checkmark	\checkmark	\checkmark	\checkmark	P	P	P	P	P	P	P	P	P	P	

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	Time-based IDS	X	X	X	\checkmark	D	_	_	_	D	_	D	_	_	D	
	ID Obfuscation	X	\checkmark	X	\checkmark	_	_	_	P	P	P	P	P	P	Р	
C	Counterattacking	X	\checkmark	\checkmark	\checkmark	P	_	_	_	P	_	P	—	_	P	
	Authentication	X	\checkmark	X	X	P	_	_	_	D	_	D	—	_	Р	
Applicati	on-level Firewall	X	\checkmark	X	X	P	_	—	—	D	—	Ρ	-	-	Р	
	CANSentry	\checkmark	\checkmark	\checkmark	\checkmark	P	P	P	Р	P	P	Р	P	P	Р	

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Conclusions

- We summarized existing DoS and spoofing attacks on CAN
- We proposed and implemented a novel stealthy selective arbitration DoS attack
- We designed a novel *CANSentry* firewall to defend against attacks that violate the CAN standard or abuse CAN's error-handling mechanism
 - *CANSentry* is the first solution that detects and prevents a broad spectrum of CAN denial and spoofing attacks
 - CANSentry does not introduce noticeable overhead or delay
 - It is very cost-effective

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