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Statistical Confidentiality Using Homomorphic Encryption

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Statistical Confidentialit y

Protection of data collected for statistical purposes.
Ensures individual privacy while enabling aggregated insights.



Statistical Dara Disclosure Control

Data Masking

Mechanisms

for Statistical

Confidentialit

- **Dara Anonymization**
- rol ity Symposium Synthetic Data
- Access Contro
- Encryption



US Census Bureau

Disclosure Avoidance 2020

"Modern computers and today's data-rich world have rendered the Census Bureau's traditional confidentiality protection methods obsolete. Those legacy methods are no match for hackers aiming to piece together the identities of the people and businesses behind published data."

"The 2020 DAS is based on a framework for assessing privacy risk known as differential privacy. It is the only solution that can respond to this threat while maximizing, the availability and utility of published census data."

https://www.census.gov/programs-surveys/decennial-census/disclosure-avoidance.2020.html#list-tab-1042747205



United States Censu s Bureau

Disclosure Avoidance (2030)





Homomorphi c Encryption

• A form of encryption that allows specific types of computations to be carried out on ciphertexts and generates an encrypted result which, when decrypted, matches the result of operations performed on the plaintexts. -Ommu

y: plaintexts

E(x + y) = E(x) + E(y)E(x * y) = E(x) * E(y)





Goal: Compute privately f(x, y) for plaintext x, y where f is a function that involves addition and multiplication (and subtraction).

Usual process:

- Encrypt the data and send it to the computing party: E(x), E(y)
- 2. Data user decrypts the data: $D(E(x)), D(E(y)) \rightarrow x, y$
- 3. Computes z = f(x, y)
- 4. Encrypts the result and sends the ciphertext to the owner t = E(z)
- 5. The owner decrypts the ciphertext to get the result D(E(t)) = z

With Homomorphic Encryption:

- 1. Encrypt the data and send it to the computing party: E(x), E(y)
- 2. Data user computes \mathcal{F} on the encrypted data and sends the ciphertext to the owner: $t = \mathcal{F}((E(x), E(y)))$
- 3. The owner decrypts the ciphertext to get the result D(t) = f(x, y)





Fully Homomorphi c Encryption over the Torus

FFHE

"Faster fully homomorphic encryption: Bootstrapping in less than 0.1 seconds", Ilaria Chillotti, Nicolas Gama, Mariya Georgieva and Malika Izabachène (Asiacrypt, 2016)

LWE Learning With Errors

https://www.zama.ai/post/estimating-the-security-of-homomorphic-schemes

https://www.zama.ai/post/tfhe-deep-dive-part-1



TFHE Ciphertexts



Concrete

- Compiler by Zama that compiles Python code into its FHE equivalent
- Implements a variant of TFHE that enables both leveled and fast bootstrapped operations as well as approximate or exact evaluation of arbitrary functions, on both booleans and integers.
- Allows writing code in Python and supports some operations of Symposium *numpy* library. Example:
 - Square
 - o Sum
 - o Min
 - Max





Limitations

- Branching (if-else conditions)
- anching () Division Negative numbers imal numbers Symposium



Computation s

• We're performing the following statistical operations: an unity symposium Mean •Median[®] Variance oMin/Max



Min/Max

Returns the minimum and maximum value of an array
We're using min() and max() methods provided by numpy library to return the min and max values



- Using sum() method provided by the numpy library to calculate the sum of all the elements of the array.
- The encrypted sum is returned, decrypted and divided by length of array to get the mean because division of encrypted numbers is not supported by concrete.

Mean



Sort the array to find the median (We implemented bubble sort)

• Used the following trick to overcome the branching limitation:

Median

- c? t f = f + c * (t f)
 Here c is the condition, t is the value to be returned if condition is true, f is the value to be returned if condition is false
- If length of array is even, we return encrypted sum of middle two numbers.
 - To find the median the sum is decrypted and divided by 2
- If length of array is odd, we simply return the middle element.



Conditionals with Ciphertexts

	if array[j] > array[j+1]: temp = array[j] array[j]=array[j+1]
	array[j+1]=temp
10	<pre>@fhe.compiler({"array": "encrypted"})</pre>
11	def sort_array(array):
12	n = array.size
13	<pre>for i in range(n):</pre>
14	for j in range(0, n - i - 1):
15	c = array[j] > array[j+1]
16	temp = array[j]
17	array[j] = array[j] + ((array[j+1] - array[j]) * c)
18	array[j + 1] = array[j + 1] + ((temp - array[j + 1]) * c)
19	return array
20	VIDE
21	<pre>inputset = [np.random.randint(0, MAX_VALUE, size=MAX_SIZE) for _ in range(50)]</pre>
22	<pre>circuit = sort_array.compile(inputset, compress_evaluation_keys=True, compress_input_ciphertexts=False)</pre>
	-Um



$V = \frac{\sum (x_i - \bar{x})^2}{n - 1}$

• Used a derived formula for variance to overcome the limitations (division). $\frac{n^2 \sum_{i=1}^{n} x_i^2 - 2n \sum_{i=1}^{n} (x_i \sum_{j=1}^{n} x_j) + n(\sum_{j=1}^{n} x_j^2)}{n^2 (n-1)}$

 The numerator is calculated in encrypted form, decrypted and divided by the denominator to get the result.

Variance









Median



• It is possible to perform some statistical operations on encrypted data.

Conclusion

- Not feasible for larger data sets especially when performing complex operations such as sorting.
 Need to develop algorithms that work with the limitations.
- Not practical for real world applications with the current resources.



Future Work

Calculate median without sort Consider other sorting algorithms i sorting Symposium

References

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[2] X. Yi, R. Paulet, and E. Bertino, "Homomorphic encryption" in Homomorphic Encryption and Applications, Springer, 2014, pp. 27– 46. [Online]. Available: <u>https://doi.org/10.1007/978-3-319-12229-</u> <u>8 2</u>

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